# Exploiting Sleep-and-Wake Strategies in the Gnutella Network

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- P2P architectures are widely used to implement large-scale collaborative networks, including file sharing systems
- Given the large sets of computing resources involved in P2P file sharing networks, their aggregate energy consumption is an important problem to be addressed
- The sleep-and-wake approach has been proposed as a general approach to reduce energy consumption in P2P systems
- Goal: evaluating how the sleep-and-wake energy-saving approach can be used to reduce energy consumption in the Gnutella network

- We introduce a general sleep-and-wake algorithm for Gnutella networks in which
  - O All leaf-peers cyclically switch between wake and sleep mode
  - O Each leaf-peer autonomously decides the time passed in sleep mode
- We define different strategies that a leaf-peer may employ to decide the duration of its sleep periods
- Such strategies have been evaluated through simulation using the general sleep-and-wake algorithm in different network scenarios

- Energy-efficient peer-to-peer systems
- Network assumptions
- General sleep-and-wake algorithm
- Sleep duration strategies
  - O VAR\_HR: duration depends on the hit rate
  - VAR\_FS: duration depends on the number of files shared
  - VAR\_QR: duration depends on the query rate
  - FIX\_nWD: duration fixed to *n* times the *wake duration*
- Performance evaluation
- Conclusions

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- Existing systems can be classified under six categories\*:
  - **Proxying:** peers can go offline to save energy by delegating some of their activities (e.g. download tasks) to proxies
  - **Task allocation optimization:** energy savings is achieved by deciding which peer will satisfy the request of another peer
  - Message reduction: energy consumption is reduced by minimizing the number of messages and the associated processing times
  - O Location-based: reduces the energy consumed by multi-hop retransmissions by improving the match between overlay and network
  - Overlay structure optimization: improves energy efficiency by controlling overlay topology or introducing new layers to the overlay
  - Sleep-and-wake: reduces energy consumption by letting peers cyclically switch between wake and sleep mode

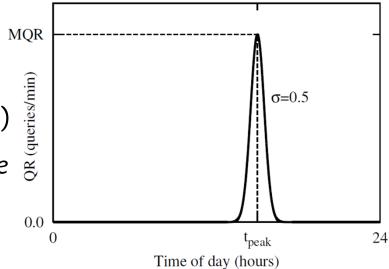
\* A. Malatras, F. Peng, B. Hirsbrunner "Energy-efficient peer-to-peer networking and overlays" in: M. S. Obaidat, A. Anpalagan, and I. Woungang (Eds.), Handbook of Green Information and Communication Systems, Elsevier, 2013

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- Two-layer overlay (Gnutella 0.6):
  - O Top layer composed of a number of ultra-peers
  - O Bottom layer comprises a higher number of leaf-peers
  - Each leaf-peer is connected to a few ultra-peers, while each ultra-peer is connected to several other ultra-peers
  - O A leaf-peer submits a query to its ultra-peers, which in turn forward the query to other ultra-peers using a TTL-limited flooding search

# Query submission rate:

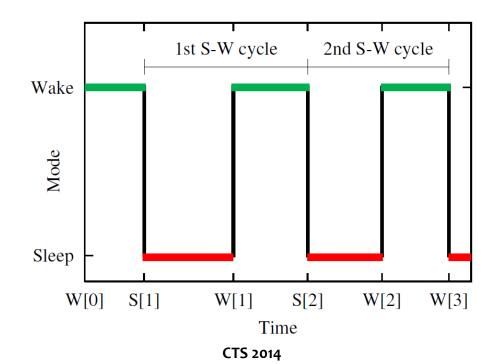
- O The inter-generation times are independent and obey an exponential distribution with a given query rate (QR)
- The QR reaches a maximum query rate (MQR) at a given time and distributes around it following a Gaussian



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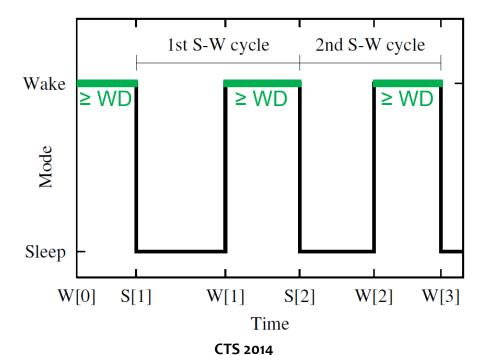
#### General sleep-and-wake algorithm (1/3)

- Leaf-peers can switch between wake and sleep mode over the time to reduce energy consumption
  - O Wake mode: the leaf-peer it is available for download requests and works at normal power level
  - Sleep mode: the leaf-peer is unavailable and works at reduced power level

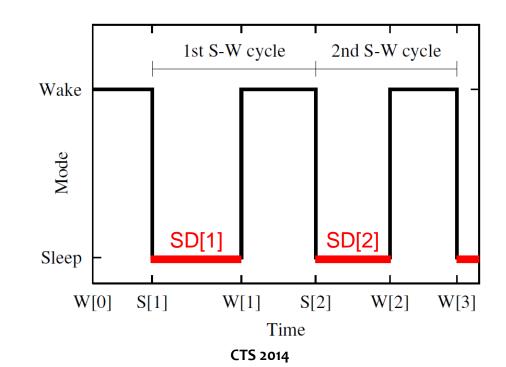


#### General sleep-and-wake algorithm (2/3)

- The duration of the i-th wake period, i.e. S[i+1]-W[i], is greater than or equal to a constant WD:
  - O It is equal to WD if at time W[i] + WD the leaf-peer is not busy with any query processing or file transfer activity
  - O Otherwise, the beginning of the next sleep period is deferred and so the i-th wake period will be longer than WD



- The duration of the i-th sleep period, SD[i], is calculated by the leaf-peer at end of the (i-1)-th wake period based on the specific strategy adopted
  - O Variable duration
  - O Fixed duration



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#### **Sleep duration strategies**

- Given the general sleep-and-wake algorithm, it is possible to define different strategies for deciding the duration of the next sleep period
- We defined and evaluated the following strategies:
  - **VAR\_HR**: variable sleep duration depending on the *hit rate*
  - VAR\_FS: variable sleep duration depending on the number of files shared
  - VAR\_QR: variable sleep duration depending on the query rate
  - **FIX\_nWD:** fixed sleep duration equal to *n* times WD

- Hit rate of the i-th wake period of a leaf-peer p, HR[i], is the number of query hits generated by p during the time interval [W[i], t] divided by t W[i], where t is the ending time of the i-th wake period
- The duration of the i-th sleep period of a leaf-peer p, denoted
  SD[i], depends on HR[i-1] as follows:

 $SD[i] = \begin{cases} 0 \ m & \text{if } HR[i-1] > 1 \ h/m \\ 5 \ m + \frac{10}{HR[i-1]} & \text{if } 0.1 \ h/m < HR[i-1] \le 1 \ h/m \\ 120 \ m & \text{otherwise} \end{cases}$ 

 Using VAR\_HR, the leaf-peers with a high hit rate will not sleep at all or will sleep for a short amount of time, while those with a lower hit rate will sleep longer With VAR\_FS, the duration of the *i*-th sleep period of a leaf-peer *p*, *SD*[*i*], depends on *FS*[*i*-1], which represents the number of files shared by *p* at the end of the (*i* - 1)-th wake period:

$$SD[i] = \begin{cases} 5 \ m & \text{if } FS[i-1] > 100 \\ 15 \ m + \frac{100 \ m}{FS[i-1]} & \text{if } 1 \le FS[i-1] \le 100 \\ 120 \ m & \text{otherwise} \end{cases}$$

 Using this strategy, the leaf-peers with a high number of files will sleep for a short amount of time, while those with a lower number of files will sleep longer

- Differently from the previous strategies, VAR\_QR links the sleep duration of a leaf-peer to its client-side behavior, i.e. the query rate of the leaf-peer during the previous wake period
- Query Rate of the *i*-th wake period of a leaf-peer *p*, denoted QR[*i*], is the number of queries submitted by *p* during the time interval [W[*i*], *t*] divided by *t* W[*i*]
- Specifically, SD[i] in VAR\_QR depends on QR[i-1] as follows:

$$SD[i] = \begin{cases} 0 \ m & \text{if } QR[i-1] > 1 \ q/m \\ 5 \ m + \frac{10}{QR[i-1]} & \text{if } 0.1 \ q/m < QR[i-1] \le 1 \ q/m \\ 120 \ m & \text{otherwise} \end{cases}$$

## FIX\_1WD and FIX\_3WD

- FIX\_1WD and FIX\_3WD are two blind strategies with which all the sleeps have the same fixed duration (introduced mostly for comparison with the previous strategies).
- Specifically, with FIX 1WD (Fixed to WD) the sleep duration is equal to WD:

$$SD[i] = WD$$

while with FIX 3WD (Fixed to 3WD), the sleep duration is equal to three times WD:

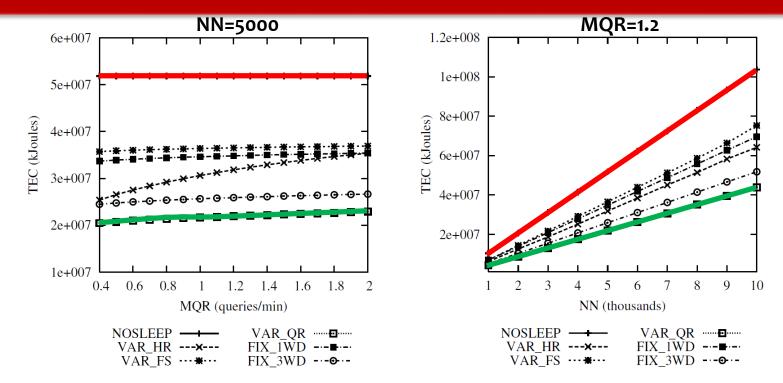
$$SD[i] = 3WD$$

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- The five strategies will be compared with a sixth strategy, referred to as NOSLEEP, in which all nodes are assumed to be always in wake mode
- Performance parameters:
  - O Total Energy Consumption (TEC) of the network
  - O Hit Rate (HR), i.e., the fraction of successful queries

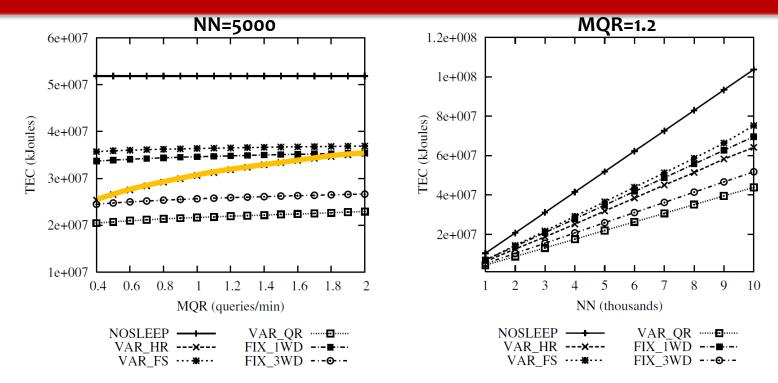
Parameter	Description	Values
NN	Number of nodes	1000-10000
MQR	Max. query rate (queries/min.)	0.4-2.0
SRD	Simulation run duration	24 h
WMPC	Wake-mode power consumption of every peer [32]	120 W
SMPC	Sleep-mode power consumption of every peer [32]	5 W
WSTT		9 s
SWTT	Sleep-to-wake transition time [33]	4 s
TTL	Time-to-live [34]	3
MTT	Message transfer time [35]	20 ms
QPT	Query processing time [36]	2 ms
QHPT	Query hit processing time [36]	1 ms
WP	Min. duration of a wake period	20 min.

# Total Energy Consumption (TEC) (1/2)



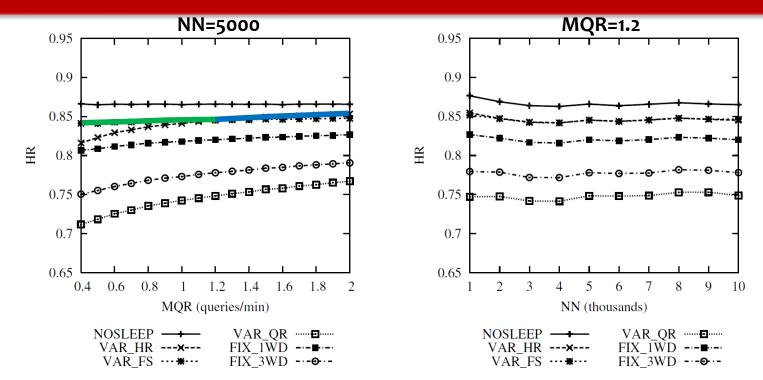
- The best result is obtained with VAR\_QR, whose TEC is on average 42% of that obtained with NOSLEEP
- TEC increases linearly with the network size → the absolute amount of energy saved increases significantly as the network grows in size

# Total Energy Consumption (TEC) (2/2)



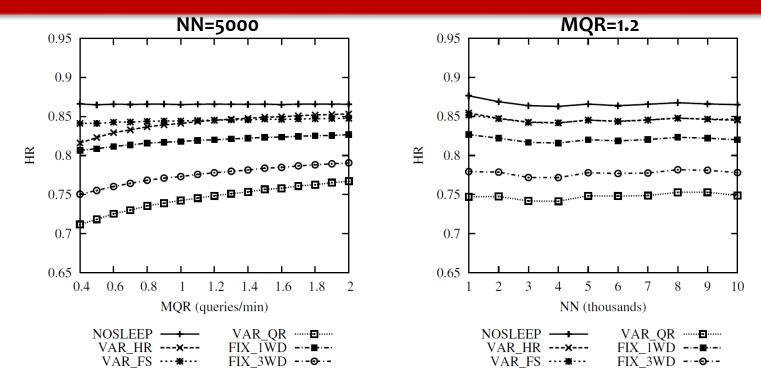
- All the strategies increase their TEC as MQR increases:
  - O High MQR values → peers submit more queries → increases the possibility that a sleep will be deferred due to the consequent client- or server-side activity
  - O More evident with VAR\_HR, because the hit rate grows proportionally with the number of queries submitted to the network

# Hit Rate (HR) (1/2)



- The highest HR is obtained with VAR\_FS and VAR\_HR: the former with MQR < 1.2, the latter with MQR > 1.2
- HR does not depend on the number of nodes: minor changes of HR are due to stochastic variations while the network is created by the simulator

# Hit Rate (HR)(2/2)



- All the strategies increase their HR as MQR increases.
  - As for the case of TEC, high MQR values increases the possibility that a sleep will be deferred due to the consequent client- or server-side activity
  - This increases the overall time passed in wake mode by the peers and consequently the possibility that a file is available when it is searched

 ESPI is an aggregate performance indicator providing an overall evaluation of a strategy X:

$$ESPI(\mathbf{X}) = \frac{\frac{TEC(\text{NOSLEEP}) - TEC(\mathbf{X})}{TEC(\text{NOSLEEP})}}{\frac{HR(\text{NOSLEEP}) - HR(\mathbf{X})}{HR(\text{NOSLEEP})}}$$

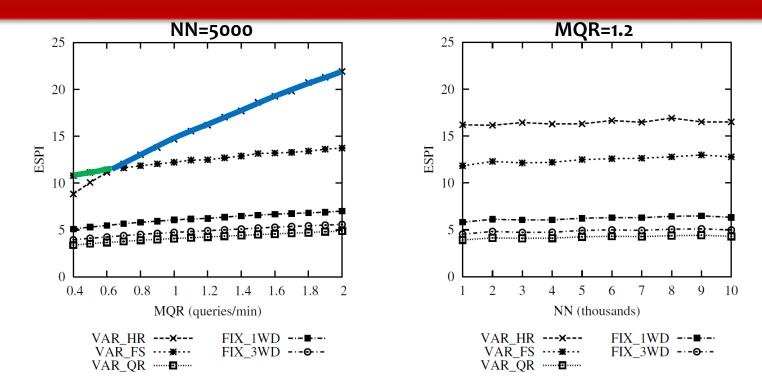
where:

- TEC(NOSLEEP): is the TEC in the NOSLEEP case
- O HR(NOSLEEP): is the HR in the NOSLEEP case
- O TEC(X): is the TEC of strategy X

O HR(X): is the HR of strategy X

O X is in {VAR\_HR, VAR\_FS, VAR\_QR, FIX\_1WD, FIX\_3WD}

#### **ESPI**



- The best strategies, based on their ESPI values, are VAR\_FS and VAR\_HR: the former with MQR < 0.8, the latter with MQR > 0.8.
- As expected, the result is independent from network size.

#### Conclusions

- Use of the sleep-and-wake energy-saving approach for reducing reduce energy consumption in Gnutella
- A general sleep-and-wake algorithm in which

O All leaf-peers cyclically switch between wake and sleep mode

O Each leaf-peer autonomously decides the time passed in sleep mode

- Different strategies that a leaf-peer may employ to decide the duration of its sleep periods have been tested.
- Simulation results have shown that the best performing strategy (VAR\_QR) consumes on average 42% of the energy consumed in a network where leaf-peers are always online